Numerical Analysis Fall 2008 MS/PhD Qualifying Examination

Instructions: Write the last four digits of your SSN (not your name) on each sheet. Complete all three problems. Clear and concise answers with good justification will improve your score.

1. ( 30 pts.) Let $A$ be an $n \times n$ real matrix and let $A^{(k)}$ denote the $k \times k$ matrix formed by the intersection of the first $k$ rows and columns of $A$. Prove that if $A^{(k)}$ is nonsingular, $k=1, \ldots, n$, then $A$ has an $L U$ factorization, where $L$ is a unit lower triangular matrix (i.e. $l_{j j}=1$ ) and $U$ is upper triangular. Express the pivots, $u_{j j}$, in terms of the determinants of the $A^{(k)}$.
2. (30 pts.) Let $A=A^{T}$ be a real $d \times d$ symmetric matrix.
i. Use Schur's Theorem to prove that $A$ has real eigenvalues and can be diagonalized by an orthogonal similarity transformation, $Q^{T} A Q=\Lambda=\operatorname{diag}\left(\lambda_{i}\right)$ where $\lambda_{i}$, $i=1, \ldots, n$ are the eigenvalues of $A$ and $Q^{T} Q=I$.
ii. Suppose $A$ has exactly $p$ positive eigenvalues. Show that there exists a $p$-dimensional subspace, $P$, of $\mathbb{R}^{d}$ such that $x^{T} A x>0$ for all $x \in P, x \neq 0$. Show that if $W$ is any subspace of dimension greater than $p$ there exists a nonzero vector $y \in W$ such that $y^{T} A y \leq 0$. (Hint: prove that there exists a nonzero vector, $y \in W$, which is orthogonal to all the eigenvectors of $A$ corresponding to positive eigenvalues.)
iii. Now consider nonsimilarity transformations, $B=S^{T} A S$ where $S$ is nonsingular. Give an example showing that the spectrum is not preserved by such transformations. Use part (ii.) to prove that $B$ has exactly as many positive, negative and zero eigenvalues as $A$.
3. (40 pts.) Let $A=A^{T}$ be a real $d \times d$ symmetric matrix. In the following we derive and analyze the MINRES method for solving $A x=b$. (Note that actual implementations of MINRES typically involve additional steps to ensure numerical stability, but here you are to consider the algorithm in exact arithmetic.)
i. Define the $k$ th Krylov subspace by:

$$
\mathcal{K}_{k}(A ; b)=\operatorname{span}\left\{b, A b, \ldots, A^{k-1} b\right\}
$$

The Lanczos process is defined by:

$$
q_{1}=\frac{b}{\|b\|_{2}}, \quad q_{0}=0, \quad \beta_{0}=0
$$

For $j=1, \ldots$

$$
\begin{gathered}
z=A q_{j}, \quad \alpha_{j}=q_{j}^{T} z, \quad z=z-\alpha_{j} q_{j}-\beta_{j-1} q_{j-1} \\
\beta_{j}=\|z\|_{2}, \quad q_{j+1}=\frac{z}{\beta_{j}}
\end{gathered}
$$

Prove that in the absence of breakdowns $\beta_{j}=0$ the vectors $\left\{q_{1}, \ldots, q_{k}\right\}$ are an orthonormal basis for the $k$ th Krylov subspace. (Hint: assuming $q_{1}, \ldots, q_{j}$ is an orthonormal basis for $\mathcal{K}_{j}(A ; b)$ project $z$ onto the orthogonal complement of $\mathcal{K}_{j}(A ; b)$ as in the Gram-Schmidt algorithm.)
ii. The $k$ th MINRES iterate, $x^{k}$, is defined as the solution of the least squares problem:

$$
\min _{y \in \mathcal{K}_{k}}\|A y-b\|_{2}
$$

Writing

$$
x^{k}=\sum_{j=1}^{k} c_{j}^{k} q_{j}
$$

Show that the expansion coefficients $c^{k}$ are solutions of the least squares problem:

$$
\min _{c \in \mathbb{R}^{k}}\left\|T^{k} c-\right\| b\left\|_{2} e_{1}^{k+1}\right\|_{2}
$$

where $T^{k}$ is the $(k+1) \times k$ tridiagonal matrix:

$$
t_{j j}=\alpha_{j}, \quad t_{j, j+1}=t_{j+1, j}=\beta_{j}
$$

and $e_{1}^{k+1}$ is the first column of the $(k+1) \times(k+1)$ identity matrix.
iii. Briefly outline (no details necessary) an efficient algorithm for computing the iterates, $c^{k}$.
iv. Show that if there are no breakdowns MINRES converges in $d$ iterations and if there is a breakdown, $\beta_{k}=0$, then $x^{k}=x$, i.e. the algorithm converges in $k$ steps.

