

28.2

As we saw, the reduction to Hessenberg form (via Givens rotations or Householder reflections) by unitary similarity transformations becomes, for a real, symmetric matrix, a reduction to ~~triangular form~~ tridiagonal form. This happens because symmetry is preserved by the unitary similarity transformation:

$$\begin{aligned} Q A Q^* &= H \implies \\ \implies (Q^*)^T A^T Q^* &= H^* - \\ \implies Q A Q^* &= H^* \implies H^* = H. \end{aligned}$$

But, naturally, a symmetric, Hessenberg matrix is tridiagonal. — o —

(b) Consider now the QR iteration for a symmetric, tridiagonal matrix.

$$\text{We have } H \equiv H_0 = Q_1 R_1$$

$$H_1 = R_1 Q_1$$

$$\text{Clearly } Q_1^* H_0 = R_1 = H_1 Q_1^*$$

$$\implies Q_1^* H_0 Q_1 = H_1$$

We already know that the QR iteration preserves the Hessenberg structure (shown in class — the book focuses entirely on symmetric matrices!). As above, since H_1 is symmetric

and Hessenberg, it must be tridiagonal.

(a) Finally following the discussion in class, which ~~discussed~~ gave the construction of the QR algorithm for general matrices (after their reduction to Hessenberg form) in terms of Givens rotations, we recall that Q is also Hessenberg (and, in fact, ~~is~~ has full upper triangle even if $\uparrow H$ is tridiagonal), while R is upper triangular.

(c) Consider the first ~~two~~ steps of the QR factorization of a tridiagonal matrix:

$$\begin{pmatrix} \cos \theta & \sin \theta & & & \\ -\sin \theta & \cos \theta & & & \\ & & I & & \end{pmatrix} \begin{pmatrix} a_{11} & a_{12} & 0 & 0 & 0 \\ a_{21} & a_{22} & a_{23} & 0 & 0 \\ 0 & a_{32} & a_{33} & a_{34} & 0 \\ 0 & 0 & * & * & * \end{pmatrix} \quad \text{where}$$

$$= \begin{pmatrix} r_{11} & r_{12} & r_{13} & 0 & 0 \\ 0 & a_{22} & a_{23} & 0 & 0 \\ 0 & a_{32} & a_{33} & a_{34} & 0 \\ 0 & 0 & * & * & * \end{pmatrix} \quad \text{where: } r_{11} = \sqrt{a_{11}^2 + a_{21}^2}, \cos \theta = \frac{a_{11}}{r_{11}}, \sin \theta = \frac{a_{21}}{r_{11}}$$

etc.

Notice that one more superdiagonal is created by the procedure!

It is easily seen that this happens at each subsequent step as well, so that R is upper triangular with only three non zero diagonals

Each step requires $\begin{cases} 10 \text{ multiplies for } R (+ 2 \text{ divs, } 2 \text{ squares,}) \\ + 5 \text{ adds} \end{cases} \begin{cases} 1 \text{ add, } 1 \text{ sqrt} \end{cases}$ to a total of $21n$ ops!

Set XII - Solutions

$$\boxed{33.1} \Rightarrow \mathcal{K}_n = \langle b, Ab, A^2b, \dots, A^{n-1}b \rangle = \langle q_1, \dots, q_n \rangle \subset \mathbb{C}^n$$

$$\text{while } K_n = [b | Ab | \dots | A^{n-1}b] = Q_n R_n.$$

Note that there is an obvious misprint in the problem statement: we must have either $x \in \mathcal{K}_n$ or $x \in \mathcal{R}(K_n)$. Either way, $\exists c_i$

$$x \in \sum_{i=1}^n c_i (A^{i-1}b) = \left(\sum_{i=1}^n c_i A^{i-1} \right) b =: p(A)b.$$

$\boxed{33.2} \Rightarrow$ The Arnoldi iteration at step n gives

$$Aq_n = h_{1n}q_1 + \dots + h_{nn}q_n + h_{n+1,n}q_{n+1} \quad (\text{eq. 33.4})$$

If $h_{n+1,n} = 0$, then the iteration "closes":

$$Aq_n = \sum_{i=1}^n h_{in}q_i \Rightarrow Aq_n \in \langle q_1, \dots, q_n \rangle$$

Equivalently, since $\mathcal{K}_n = \langle b, Ab, \dots, A^{n-1}b \rangle = \langle q_1, \dots, q_n \rangle$

then $A^n b = A \left(\sum_{i=1}^n c_i q_i \right) = \sum_{i=1}^{n-1} c_i (Aq_i) + c_n Aq_n \in \mathcal{K}_n$,

and $\mathcal{K}_{n+1} = \langle b, \dots, A^{n-1}b, A^n b \rangle = \langle b, \dots, A^{n-1}b \rangle = \mathcal{K}_n$.

(a) When $h_{n+1,n} = 0$, equation (33.13) simplifies to

$$AQ_n = Q_n H_n \quad (H_n \text{ as defined in eq. 33.9})$$

If the full Hessenberg form contains a zero

subdiagonal entry, $h_{n+1,n} = 0$, then the

eigenvalue problem for the matrix A is reduced

and $A^{n+1}b \in \mathcal{K}_n \Rightarrow \mathcal{K}_{n+2} = \langle b, \dots, A^n b, A^{n+1}b \rangle = \langle b, A^{n+1}b \rangle = \mathcal{K}_n$
 and the same is true for any \mathcal{K}_ℓ , $\ell > n$.

(d) This was already shown in (a), as part of the discussion of the structure of the Hessenberg reduction.

(e) If A is nonsingular then $\{A^i b\}_{i=0}^{n-1}$ are
~~(*) independent: $\sum_{i=1}^n c_i A^i b = 0 \Rightarrow$~~ the n vectors

$\{Ab, A(Ab), \dots, A(A^{n-1}b)\}$ are independent: indeed
 $\sum_{i=1}^n c_i A^i b = A \left\{ \sum_{i=1}^n c_i A^{i-1} b \right\} = 0 \stackrel{(A \text{ non-singular})}{\Rightarrow} \sum_{i=1}^n c_i A^{i-1} b = 0$

But then $c_i = 0$, $i=1, \dots, n$ since the vectors $\{A^{i-1}b\}_{i=1}^n$ are independent. Since \mathcal{K}_n invariant, the vectors $\{A^i b\}_{i=0}^{n-1}$ are all in \mathcal{K}_n , and thus they form a basis.

Let $b = \sum_{i=1}^n d_i A^i b \Rightarrow Ax = \sum_{i=1}^n d_i A^i b$
 $\Rightarrow x = \sum_{i=1}^n d_i A^{i-1} b \in \mathcal{K}_n$.

34.1 $A \in \mathbb{C}^{m \times m}$, $b \in \mathbb{C}^m$, $p \in \mathbb{P}^n$ (monic polys., highest coeff. = 1)
Computing $p(A)b$

$$p(z) = c_0 + z(c_1 + \dots + z(c_{m-1} + z) \dots)$$

(a)
$$P = A + c_{m-1}I$$

for $i = 2:m$
$$P = AP + c_{m-i}I$$

end
$$\rightarrow p(A) = c_0 I + A(c_1 I + \dots + A(c_{m-1} I + A))$$

(Note: the summation only touches diagonal entries)

$$\text{Total cost: } 2(m-1) * m^2 + 2m(m-1) = \Theta(m^4)$$

(b) Now: $pb = Ab + c_{m-1}b$

for $i = 2:m$

$$pb = A * pb + c_{m-i}b$$

end

$$\text{Cost: } (m-1) * 2m^2 + 2m(m-1) = \Theta(m^3) (\sim 2m^3)$$

(c) "Obvious method": (i) need $A^k b$, $k = 1, \dots, m-1$

Total cost: $(m-1) * 2m^2$ (successively compute terms)

(ii) $c_k A^k b$: $m * m = m^2$

(iii) Additions: $m * m = m^2$

$$\text{Total: } \Theta(m^3) (\sim 2m^3)$$

i.e. (b)-(c) cost the same!

35.4 To minimize:

$\|r_n\| = \|H_n y - \|b\|e_1\|$: need QR factorization of H_n . This is accomplished by $(n-1)$ Givens rotations, as follows:

$$\begin{pmatrix} (c_1 & -s_1) \\ & I_{n-2} \\ (s_1 & c_1) \end{pmatrix} \begin{pmatrix} h_{11} & h_{12} & \dots & h_{1n} \\ h_{21} & h_{22} & \dots & h_{2n} \\ & h_{32} & \dots & \\ & & \dots & \\ & & & h_{n-1,n} & h_{nn} \end{pmatrix}$$

c_1, s_1 chosen to eliminate h_{21} :

$$h'_{21} = s_1 h_{11} + c_1 h_{21} = 0 \Rightarrow s_1 = \frac{h_{21}}{\sqrt{h_{11}^2 + h_{21}^2}}, \quad c_1 = -\frac{h_{11}}{\sqrt{h_{11}^2 + h_{21}^2}}$$

Proceeding similarly, we arrive at $H_n = Q_n R_n$

where $Q_n = G_{n-1}^T \dots G_1^T$

Each step $\square \mathcal{O}(n)$ (total n^2); then

$$Q_n^* (H_n y = \|b\|e_1) \Rightarrow R_n y = \underbrace{(Q_n^* e_1)}_{\substack{\text{just steal } 1^{\text{st}} \text{ column} \\ \mathcal{O}(n)}}} \|b\|$$

Solve by backsubstitution in $\mathcal{O}(n^2)$.

Going to H_{n+1} , can reuse all the same G_i ;
 Additional work: $(G_{n-1} \dots G_1) h_{(n+1)}$ ($n+1$ column) is $\mathcal{O}(n)$
 Eliminating h_{n+1} requires one additional G , G_n ($\mathcal{O}(n)$).